

# Rainbow matchings for 3-uniform hypergraphs

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# Erdős Matching Conjecture

## Erdős Matching Conjecture, 1965

For positive integers  $k, n, s$ , if  $H$  is a  $k$ -graph on  $n$  vertices and  $\nu(H) < s$ , then

$$e(H) \leq \max \left\{ \binom{ks-1}{k}, \binom{n}{k} - \binom{n-s+1}{k} \right\}.$$

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## Theorem (Erdős, 1965)

*The conjecture is true for  $n > n_0(k, s)$ .*

# Erdős Matching Conjecture

- Frankl, 2013

The conjecture is true for  $n > (2s - 1)k - s$ .

- Frankl and Kupavskii, 2018

The conjecture is true for  $s \geq s_0$  and  
 $n \geq \frac{5}{3}(s - 1)k - \frac{2}{3}(s - 1)$ .

For 3-graphs:

- Frankl, 2017

The conjecture is true for  $k = 3$ .

# Rainbow Matching

Conjecture (Huang, Loh and Sudakov, 2012; independently by Aharoni and Howard)

Let  $F_1, \dots, F_t$  be  $k$ -graphs on  $[n]$ . If

$$|F_i| > \max \left\{ \binom{n}{k} - \binom{n-t+1}{k}, \binom{kt-1}{k} \right\}$$

for all  $1 \leq i \leq t$ , then there is a 'rainbow' matching of size  $t$ : one that contains exactly one edge from each family.

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Theorem (Huang, Loh and Sudakov, 2012)

*The conjecture above is true for  $n > 3k^2t$ .*

# Rainbow Matching

## Conjecture

Let  $F_1, \dots, F_t$  be  $k$ -graphs on  $[n]$ . If

$$\delta_1(F_i) > \binom{n-1}{k-1} - \binom{n-t}{k-1}$$

for all  $1 \leq i \leq t$ , then there is a 'rainbow' matching of size  $t$ .

## Definition

$\delta_1(H)$ : minimum vertex degree in  $H$

## Dirac Type Problem

### Theorem (Dirac, 1952)

*A simple graph  $G$  with  $n$  vertices ( $n \geq 3$ ) is Hamiltonian if  $\delta(G) \geq n/2$ . In particular, if  $n$  is even, then  $G$  contains a perfect matching.*

### Theorem (Kühn, Osthus and Treglown, 2013; independently, Khan, 2013)

*There exists an  $n_0 \in \mathbb{N}$  such that the following holds. Suppose that  $H$  is a 3-uniform hypergraph whose order  $n > n_0$  is divisible by 3. If*

$$\delta_1(H) > \binom{n-1}{2} - \binom{2n/3}{2}$$

*then  $H$  has a perfect matching.*



## Our Result

### Theorem (Lu, Yu, and Y., 2021)

Let  $n \in 3\mathbb{Z}$  be sufficiently large, and let  $\mathcal{F} = \{F_1, \dots, F_{n/3}\}$  be a family of 3-graphs such that  $|V(F_i)| = n$  and  $V(F_i) = V(F_1)$  for  $i \in [n/3]$ . If

$$\delta_1(F_i) > \binom{n-1}{2} - \binom{2n/3}{2}$$

for  $i \in [n/3]$ , then  $\mathcal{F}$  admits a rainbow matching.

## Proof Idea

- We convert it to the problem of finding a perfect matching in a balanced  $(1, 3)$ -partite 4-graph.
- For any integer  $k \geq 3$ , a  $k$ -graph  $H$  is  $(1, k - 1)$ -partite if there exists a partition of  $V(H)$  into sets  $V_1, V_2$  (called *partition classes*) such that for any  $e \in E(H)$ ,  $|e \cap V_1| = 1$  and  $|e \cap V_2| = k - 1$ .
- A  $(1, k - 1)$ -partite  $k$ -graph with partition classes  $V_1, V_2$  is **balanced** if  $(k - 1)|V_1| = |V_2|$ .

## Proof Idea

- We define a  $(1, 3)$ -partite 4-graph  $H$  with respect to  $\mathcal{F}$  by letting
  - $V(H) = P \cup Q$ ,  
where  $P = V(F_i)$  and  $Q = \{v_1, \dots, v_{n/3}\}$ ,  
and
  - $N_H(v_i) = F_i$ , for each  $i \in [n/3]$ .
- Then we have

$$\delta_1(N_H(v_i)) > \binom{n-1}{2} - \binom{2n/3}{2}$$

for each  $i \in [n/3]$ .

## Proof Idea

For the case when  $H$  is close to  $H_0$ , we prove the conjecture is true for  $k = 3$  and all  $t \in [n/3]$  (by induction).

- Base:  $t \leq n/200$ , greedy construction
- $t \geq n/400$ 
  - ◇  $H$  is close to  $H_0$  at every vertex
  - ◇  $H$  is not close to  $H_0$  at some vertices


## Proof Idea

When  $H$  is not close to  $H_0$ , we follow an approach by Alon, Frankl, Huang, Rödl, Ruciński, and Sudakov.

- Find an absorber  $M_{abs}$
- Find an almost perfect matching  $M'$  in  $H \setminus V(M_{abs})$ 
  - ◇ Find perfect fractional matchings in random subgraphs of  $H - V(M_{abs})$
  - ◇ Convert to an almost perfect matching
- Use  $M_{abs}$  to extend  $M'$  to a perfect matching

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# Absorbing Lemma

## Lemma


Let  $n \in 3\mathbb{Z}$  be large enough and let  $H$  be a  $(1, 3)$ -partite 4-graph with partition classes  $Q, P$  such that  $3|Q| = |P|$  and

$$\delta_1(H) \geq \frac{n}{3} \left( \binom{n-1}{2} - \binom{2n/3}{2} + 1 \right).$$

Let  $\rho, \rho'$  be constants such that  $0 < \rho' \ll \rho \ll 1$ . Then  $H$  has a matching  $M'$  such that  $|M'| \leq \rho n$  and, for any subset  $S \subseteq V(H)$  with  $|S| \leq \rho' n$  and  $3|S \cap Q| = |S \cap R|$ ,  $H[S \cup V(M)]$  has a perfect matching.

## Proof Idea

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## Almost Perfect Matching

- Form a random subgraph  $R \subseteq H$  by taking each vertex with probability  $n^{-0.9}$ . We take  $n^{1.1}$  independent copies of  $R$ .
- With high probability, all those copies have certain properties, for example, containing a perfect fractional matching.
- Those properties enable us to perform another round of random sampling to find a spanning subgraph satisfying the conditions of a 'Rödl Nibble' theorem .

## Proof Idea

When  $H$  is not close to  $H_0$ , we follow an approach by Alon, Frankl, Huang, Rödl, Ruciński, and Sudakov.

- Find an absorber  $M_{abs}$
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# Fractional Matching

## Lemma

Let  $\rho, \varepsilon$  be constants with  $0 < \varepsilon \ll 1$  and  $0 < \rho < \varepsilon^{12}$ , and let  $H$  be a  $(1, 3)$ -partite 4-graph with partition classes  $Q, P$  (with  $3|Q| = |P|$ ). Suppose

$$d_H(\{u, v\}) > \binom{n-1}{2} - \binom{2n/3}{2} - \rho n^2$$

for any  $v \in Q$  and  $u \in P$ . If  $H$  contains no independent set  $S$  with  $|S \cap Q| \geq n/3 - \varepsilon^2 n$  and  $|S \cap P| \geq 2n/3 - \varepsilon^2 n$ , then  $H$  contains a perfect fractional matching.

# Fractional Matching

- We bound the size of independent sets in each random induced subgraph using a hypergraph container result. (“not close to the extremal graph”)
- Using the Strong Duality Theorem, we would like to convert this problem to finding a perfect matching in a stable family, and deal with it by Tutte-Berge formula.

# Fractional Matching

- We bound the size of independent sets in each random induced subgraph using a hypergraph container result. (“not close to the extremal graph”)
- Using the Strong Duality Theorem, we would like to convert this problem to finding a perfect matching in a stable family, and deal with it by Tutte-Berge formula.
- So we need to take the first round of random sampling more carefully such that each induced subgraph taken is balanced.

## Balanced Induced Subgraphs

### Lemma

Let  $S \subset V(H)$  be a set of vertices such that

$$|S \cap Q| = n^{0.99}/3 \quad \text{and} \quad |S \cap P| = n^{0.99}.$$


Let  $R_+^i$  be chosen from  $V(H)$  by taking each vertex uniformly at random with probability  $n^{-0.9}$ , for each  $i \in [n^{1.1}]$ , independently.

Define  $R_-^i = R_+^i \cap (V(H) \setminus S)$ ,  $1 \leq i \leq n^{1.1}$ .

Then, with probability  $1 - o(1)$ , there exist  $R_i$ ,  $i \in [n^{1.1}]$ , such that  $R_-^i \subseteq R^i \subseteq R_+^i$  and  $R^i$  is balanced.

## Proof Idea

When  $H$  is not close to  $H_0$ , we follow an approach by Alon, Frankl, Huang, Rödl, Ruciński, and Sudakov.

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# Almost Perfect Matching

## Theorem (Pippenger and Spencer, 1989)

*For every integer  $k \geq 2$  and real  $r \geq 1$  and  $a > 0$ , there are  $\gamma = \gamma(k, r, a) > 0$  and  $d_0 = d_0(k, r, a)$  such that for every  $n$  and  $D \geq d_0$  the following holds: Every  $k$ -uniform hypergraph  $H = (V, E)$  on a set  $V$  of  $n$  vertices in which all vertices have positive degrees and which satisfies the following conditions:*

- (1) For all vertices  $x \in V$  but at most  $\gamma n$  of them,  $d(x) = (1 \pm \gamma)D$ ;*
- (2) For all  $x \in V$ ,  $d(x) < rD$ ;*
- (3) For any two distinct  $x, y \in V$ ,  $d(x, y) < \gamma D$ ;*

*contains a cover of at most  $(1 + a)(n/k)$  edges.*



# Almost Perfect Matching

## Lemma

Let  $\sigma > 0$  and  $0 < \rho \leq \varepsilon/3 \ll 1$ . Let  $H$  be a  $(1, 3)$ -partite 4-graph with partition classes  $Q, P$ , where  $|Q| = n/3$  and  $|P| = n$ .

Suppose  $H$  is not  $\varepsilon$ -close to  $H_{1,3}(n, n/3)$  and

$$d_H(\{u, v\}) \geq \binom{n-1}{2} - \binom{2n/3}{2} - \rho n^2$$

for any  $v \in Q$  and  $u \in P$ . Then  $H$  contains a matching covering all but at most  $\sigma n$  vertices.

## Proof Idea

When  $H$  is not close to  $H_0$ , we follow an approach by Alon, Frankl, Huang, Rödl, Ruciński, and Sudakov.

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## Proof Idea

Thus the balanced  $(1, 3)$ -partite 4-graph  $H$  defined by  $\mathcal{F}$  has a perfect matching, and hence the following holds.

### Theorem (Lu, Yu, and Y., 2021)

Let  $n \in 3\mathbb{Z}$  be sufficiently large, and let  $\mathcal{F} = \{F_1, \dots, F_{n/3}\}$  be a family of 3-graphs such that  $|V(F_i)| = n$  and  $V(F_i) = V(F_1)$  for  $i \in [n/3]$ . If

$$\delta_1(F_i) > \binom{n-1}{2} - \binom{2n/3}{2}$$

for  $i \in [n/3]$ , then  $\mathcal{F}$  admits a rainbow matching.

# Thank you!

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