

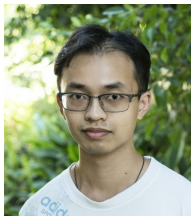
Uniqueness and Rapid Mixing in the Bipartite Hardcore Model

Xiaoyu Chen

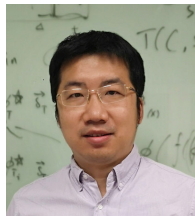


Nanjing University

based on joint work with



Jingcheng Liu



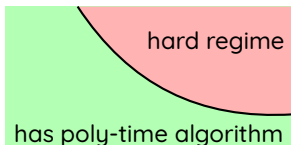
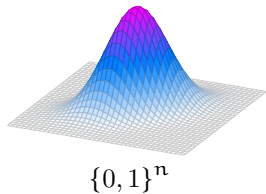
Yitong Yin

Sampling problem:

Draw (approximate) random samples from a distribution

Gibbs distribution:

- ▶ high-dimensional joint distribution
- ▶ described by few parameters and local interactions
- ▶ typical input size: $\text{poly}(n)$



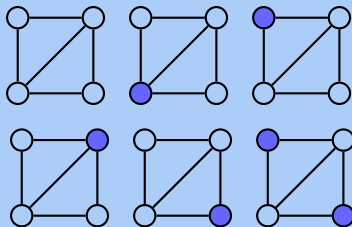
Computational phase transition:

computational complexity of sampling problem changes sharply around certain parameter values

Counting & sampling independent set (#IS)

- ▶ $G = ([n], E)$ with n vertices and max degree Δ .
- ▶ #IS: how many independent sets are there in G ?

An example



number of
independent set: 6

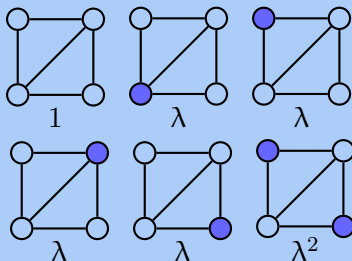
- ▶ exact counting independent set is #P-hard.
- ▶ approximate counting $\stackrel{[JVV86]}{\iff}$ approximate sampling:
 - ▶ $\Delta \leq 5$, poly-time algorithm for approx. sampling [Wei06]
 - ▶ $\Delta \geq 6$, no poly-time algorithm unless $\mathbf{NP} = \mathbf{RP}$ [Sly10]

Hardcore model

- ▶ $G = ([n], E)$ with n vertices and max degree Δ .
- ▶ Fugacity $\lambda > 0$ is a real number.
- ▶ $\text{Ind}(G) = \{S \subseteq [n] \mid S \text{ is an independent set}\}$.
- ▶ Gibbs distribution

$$\forall S \in \text{Ind}(G), \quad \mu(S) := \frac{\lambda^{|S|}}{Z}, \quad \text{where } Z_G(\lambda) := \sum_{I \in \text{Ind}(G)} \lambda^{|I|}.$$

An example



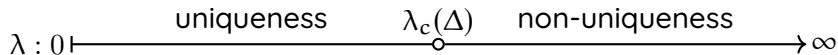
Partition function:

$$Z = 1 + 4\lambda + \lambda^2$$

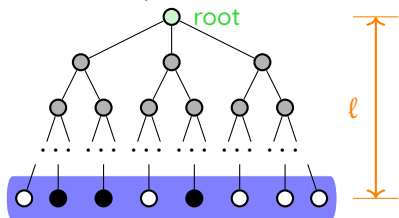
This model is self-reducible

Computational phase transition

On Δ -regular tree:



tree uniqueness threshold: $\lambda_c(\Delta) := (\Delta - 1)^{(\Delta-1)} / (\Delta - 2)^\Delta \approx \frac{e}{\Delta}$



Tree uniqueness

$\Pr_{S \sim \mu} [\text{root} \in S \mid \sigma]$ does not depend on σ when $\ell \rightarrow \infty$

if and only if $\lambda \leq \lambda_c(\Delta)$

σ : boundary condition on level ℓ

On general graph with maximum degree Δ :

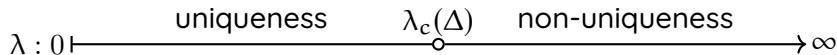


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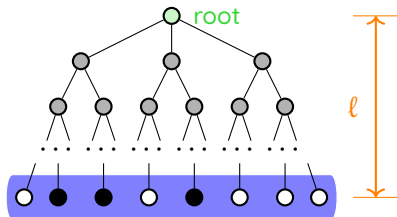
- ▶ $\lambda < \lambda_c$: poly-time algorithm for approx. sampling [Wei06]
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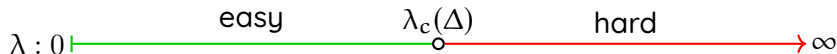
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Hardcore model on bipartite graph (weighted #BIS)

It is easy: there is a poly-time algorithm to find a maximum independent set in the bipartite graph (König's theorem¹).

It is hard: many important problems are proved to be #BIS-equivalent or #BIS-hard under AP-reductions.

Selected examples

- ▶ stable matchings (counting)
- ▶ ferro. Potts model (parti. func.)
- ▶ ferro. Ising with mixed external fields (parti. func.)

[DGGJ04, GJ07, DGJ10, CGM12 DGJR12, GJ12a, BDG+13, LLZ14, GJ15, CGG+16, GŠVY16, GGY21]

Conjecture[DGGJ04]:

#BIS represents an intermediate complexity class:

- ▶ it has no FPRAS in general
- ▶ it is easier than #SAT

¹In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.

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Previous algorithmic results

Non-uniqueness regime:

- ▶ α -expander bipartite graph:
 - ▶ $\lambda \geq (C_0\Delta)^{4/\alpha}$, an $n^{O(\log \Delta)}$ time sampler [JKP20]
 - ▶ $\lambda \geq (C_1\Delta)^{6/\alpha}$, an $O(n \log n)$ time sampler [CGG+21]
 - ▶ $\lambda \geq (C_2\Delta)^{2/\alpha}$, an $n^{O(\log \Delta)}$ time sampler [FGKP23]
- ▶ Δ -regular α -expander bipartite graph:
 - ▶ $\lambda \geq \frac{f(\alpha) \log \Delta}{\Delta^{1/4}}$, an $n^{O(\Delta)}$ time sampler [JPP22]
- ▶ random Δ -regular bipartite graph:
 - ▶ $\Delta \geq \Delta_0$, $\lambda \geq \frac{\log^4 \Delta}{\Delta}$, an $n^{O(1)}$ time sampler [LLLLM19]
 - ▶ $\Delta \geq \Delta_1$, $\lambda \geq \frac{50 \log^2 \Delta}{\Delta}$, an $n^{1+O(\frac{\log^2(\Delta)}{\Delta})}$ time sampler [JKP20]
 - ▶ $\Delta \geq \Delta_2$, $\lambda \geq \frac{100 \log \Delta}{\Delta}$, an $O(n \log n)$ time sampler [CGŠV22]
- ▶ unbalanced bipartite graph:
 - ▶ $6\Delta_L \Delta_R \lambda \leq (1 + \lambda)^{\frac{\delta_R}{\Delta_L}}$, an $n^{O(\log(\Delta_L \Delta_R))}$ time sampler [CP20]
 - ▶ $3.4\Delta_L \Delta_R \lambda \leq (1 + \lambda)^{\frac{\delta_R}{\Delta_L}}$, an $n^{O(\log(\Delta_L \Delta_R))}$ time sampler [FGKP23]
 - ▶ $(1 + e)\Delta_L \Delta_R \lambda \leq (1 + \lambda)^{\frac{\delta_R}{\Delta_L}}$, an $O(n \log n)$ time sampler [BCP22]

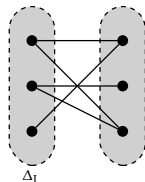
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Uniqueness regime:

[... , AJKPV22, CFYZ22, CE22]

- ▶ general graph: if $\lambda < \lambda_c(\Delta)$, there is an $O(n \log n)$ time sampler
- ▶ bipartite graph: if $\lambda = 1$, $\Delta_L \leq 5$, an $O(n^2)$ time sampler [LL15]
($\lambda = 1 \wedge \lambda < \lambda_c(\Delta) \Leftrightarrow \Delta \leq 5$)

$$\lambda_c(\Delta) = \frac{(\Delta - 1)^{\Delta - 1}}{(\Delta - 2)^\Delta}$$



Our results

For $\delta \in (0, 1)$, $\Delta_L \geq 3$, if $\lambda \leq (1 - \delta)\lambda_c(\Delta_L)$, then

- ▶ the system is in the uniqueness regime
- ▶ there is a sampler that runs in time

$$T := n \left(\frac{\Delta_L \log n}{\lambda} \right)^{O(1/\delta)}$$

- ▶ the mixing time of Glauber dynamics is bounded by $O(n^2) \cdot T$

- ▶ When $\Delta_L = 1$, G is a forest, which is trivial.
- ▶ When $\Delta_L = 2$, this model becomes an Ising model. Our results still work, but since $\lambda_c(2) = \infty$, it is quite technical to state them here.

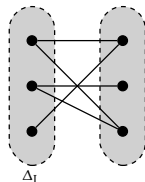
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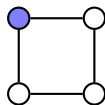
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Glauber dynamics for Hardcore model:

start from an arbitrary independent set X_0 ;

for t from 1 to T **do**:

- ▶ pick a vertex $v \in V$ uniformly at random;
- ▶ with prob. $\frac{\lambda}{1+\lambda}$, let $S = X_{t-1} \cup \{v\}$;
with prob. $\frac{1}{1+\lambda}$, let $S = X_{t-1} \setminus \{v\}$;
- ▶ **if** $S \in \text{Ind}(G)$ **then** $X_t = S$ **else** $X_t = X_{t-1}$;



irreducible + aperiodic + reversible $\implies X_t \sim \mu$ as $t \rightarrow \infty$

mixing time: essential running time of Glauber dynamics

$$T_{\text{mix}} := \max_{X_0} \min\{t \mid D_{\text{TV}}(X_t \parallel \mu) \leq 1/100\}$$

total variation distance: cononical distance between distributions

$$D_{\text{TV}}(X_t \parallel \mu) := \frac{1}{2} \sum_{S \in \text{Ind}(G)} |\Pr[X_t = S] - \mu(S)|$$

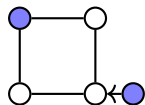
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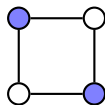
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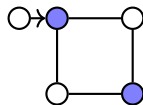
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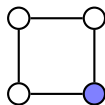
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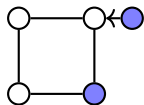
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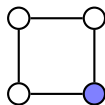
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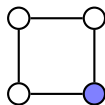
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Background

Proof outline

Fast sampler

Mixing of Glauber dynamics on $L \cup R$

Background

Let ν be a distribution over $\Omega = \{-1, +1\}^n$. $\forall \sigma \in \Omega$, $\|\sigma\|_+ = |\{i \mid \sigma_i = 1\}|$

impose external field $\theta > 0$

$\theta * \nu$: a distribution on Ω :

$$\forall \sigma, \quad (\theta * \nu)(\sigma) \propto \nu(\sigma) \cdot \theta^{\|\sigma\|_+}$$

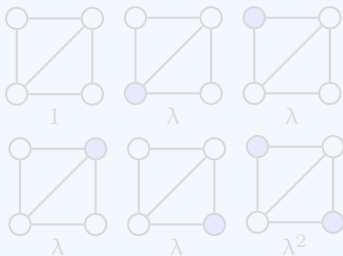
flip the distribution

$\bar{\nu}$: a distribution on Ω :

$$\forall \sigma, \quad \bar{\nu}(\sigma) = \nu(-\sigma)$$

► hardcore model: μ (fugacity λ) $\implies \theta * \mu$ (fugacity $\theta\lambda$)

An example



μ : hardcore model with fugacity λ

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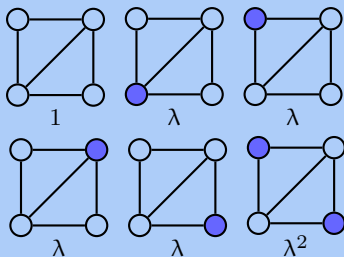
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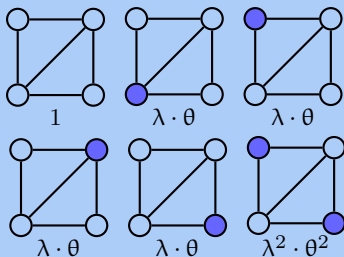
flip the distribution

$\bar{\nu}$: a distribution on Ω :

$$\forall \sigma, \quad \bar{\nu}(\sigma) = \nu(-\sigma)$$

► hardcore model: μ (fugacity λ) $\implies \theta * \mu$ (fugacity $\theta\lambda$)

An example



$\theta * \mu$: hardcore model with fugacity $\theta\lambda$

Background

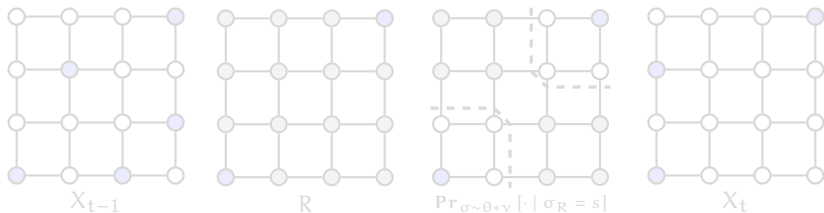
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► an example: $\nu = \mu$ is a hardcore distribution and $\theta \in (0, 1)$, $s = +1$.



irreducible + aperiodic + reversible [CFYZ21] $\implies X_t \sim \nu$ as $t \rightarrow \infty$ 🤔

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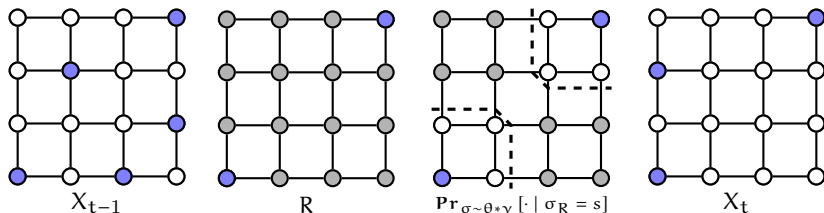
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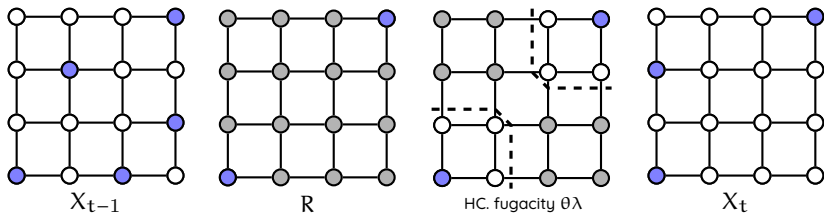
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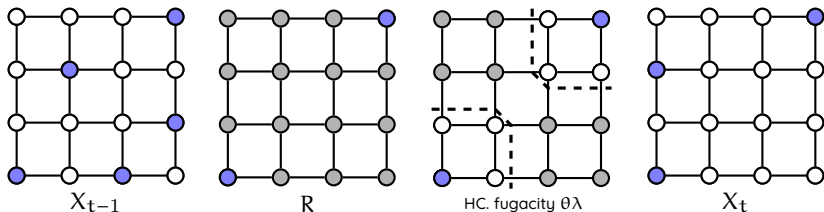
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$$\Psi_\nu(i, j) := \Pr[X_j = +1 \mid X_i = +1] - \Pr[X_j = +1 \mid X_i = -1] = \frac{\text{Cov}(X_i, X_j)}{\text{Var}(X_i)}$$

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$$\lambda_{\max}(\Psi_\nu) \leq \eta \iff \|\Psi_\nu\|_\infty \leq \eta$$

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Ψ_ν is similar to $\text{Corr}(X)$

K-marginal stability [CFYZ22, CE22]

there is $\rho \in \{\nu, \bar{\nu}\}$ that for $i \in [n]$, $S \subseteq \Lambda \subseteq [n] \setminus \{i\}$, $\tau \in \Omega(\rho_\Lambda)$,

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Some requirement on the marginal probability of ν .

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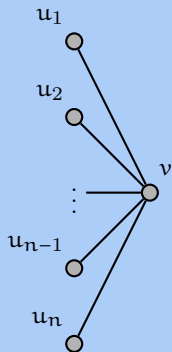
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Proof outline

example



Let $\lambda = 1$ be the fugacity

μ : Gibbs distribution of the hardcore model

$\|\Psi_\mu\|_\infty$ is unbounded

▶ $\|\Psi_\mu\|_\infty = \frac{n}{2}$

What could we do? 🤔

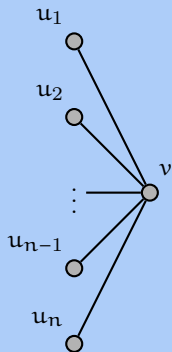
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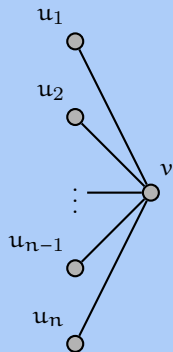
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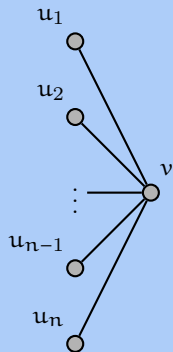
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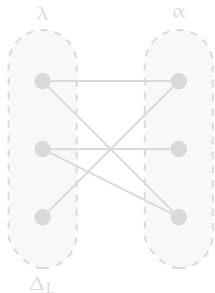
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Glauber dynamics mixes in $\tilde{O}(n)$

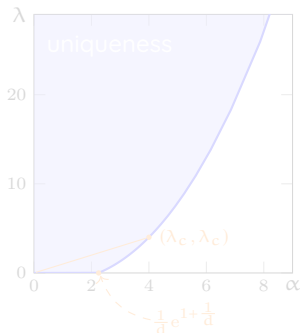
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Let $d = \Delta_L - 1$, this parametric curve is the boundary of our uniqueness regime, for $w > d^{-1}$:

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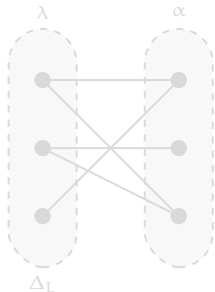
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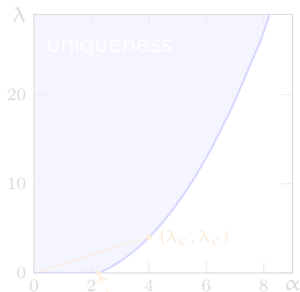
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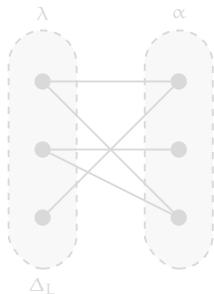
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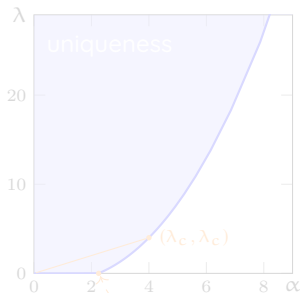
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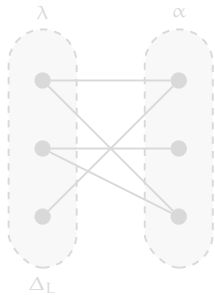
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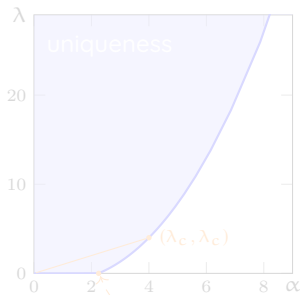
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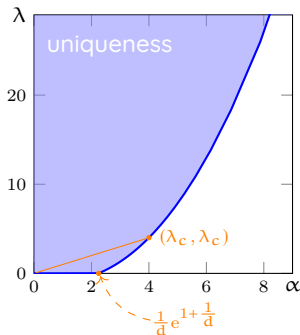
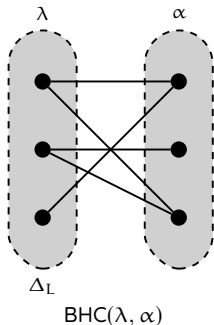
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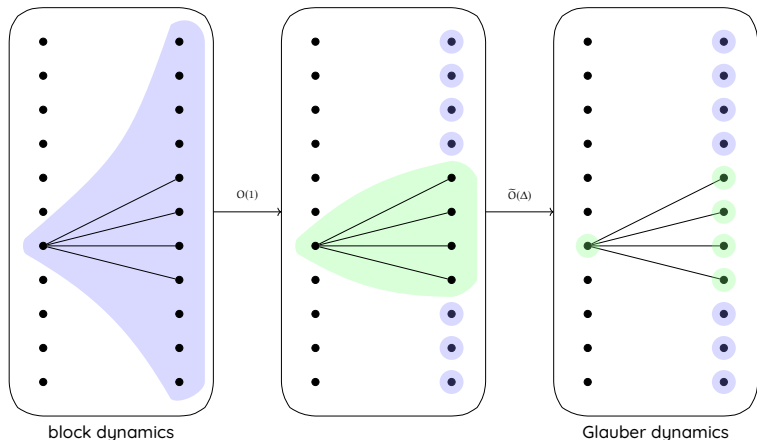
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Proof outline: mixing of GD on μ

- ▶ Glauber dynamics on $\nu = \mu_L$ is rapidly mixing.
- ▶ It works like a block dynamics that update a random vertex on the left and all the vertices on the right in each step.
- ▶ We finish the proof by comparing it and the Glauber dynamics on μ via the **block factorization** [CMT15, CP20, CLV21].



Thank you

arXiv:2305.00186

Summary

For $\delta \in (0, 1)$, $\Delta_L \geq 3$, if $\lambda \leq (1 - \delta)\lambda_c(\Delta_L)$, then

- ▶ the system is in the uniqueness regime
- ▶ there is a sampler that runs in time

$$T := n \left(\frac{\Delta_L \log n}{\lambda} \right)^{O(1/\delta)}$$

- ▶ the mixing time of Glauber dynamics is bounded by $O(n^2) \cdot T$

Open problems

- ▶ Remove the dependency on Δ_L in the running time of the sampler.
- ▶ Better mixing time for the Glauber dynamics.
- ▶ Bipartite hardcore model for negative λ .