

# On the polynomial reconstruction of digraphs and graphs

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- Edge Reconstruction Conjecture of Graphs
- Reconstruction of Characteristic Polynomials

## 2 Edge Reconstruction of Digraph Polynomials

- Two Equalities on Determinants and Permanents
- Six Digraph Polynomials
- Main Results

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# Vertex Reconstruction Conjecture of Graphs

- The famous **Ulam's conjecture** states that each simple graph  $G = (V, E)$  with  $n \geq 3$  vertices be uniquely reconstructed from its **deck**  $\{G - v | v \in V(G)\}$ . Each  $G - v$  is called a **card** of  $G$ .

Obviously,  $G \Rightarrow \{G - v | v \in V\}$ .

Ulam's Conjecture :  $\{G - v | v \in V\} \Rightarrow \text{unique } G \text{ if } n \geq 3?$

- According to reliable sources (Kelly's doctoral thesis appeared in 1942), it was discovered in Wisconsin in 1941 by Kelly and Ulam .



S. M. Ulam, *A Collection of Mathematical Problems*, Wiley (Interscience), New York, 1960, p29.



J. A. Bondy, R. L. Hemminger, *Graph Reconstruction—Survey*, Journal of Graph Theory, 1(1977), 227-268.

- The trees are reconstructible (Manvel, *Canad. J. Math.*, 1970)).
- The maximal outerplanar graphs are reconstructible (Manvel, *DM*, 1972).
- Regular graphs are reconstructible (The proof is easy).
- The unicyclic graphs are reconstructible (Arjomandi, Corneil, *Canad. J. Math.*, 1974).
- The disconnected graphs are reconstructible (Manvel, *JCTB*, 1976).

# Vertex Reconstruction Conjecture of Graphs

- McKay verified by computer that the conjecture holds for  $3 \leq V(G) \leq 10$ . (McKay, JGT, 1977).
- Godsil and McKay (JCTB, 1981) proved that a graph is reconstructible if all but at most one of its eigenvalues are simple and have eigenvectors not orthogonal to the vector  $j$  with all entries equal to one.
- Hong (JCTB, 1982) proved that if there exists a card  $G - v$  of  $G$  none of whose eigenvectors is orthogonal to the vector  $j$  with all entries equal to one, then  $G$  is reconstructible from its deck.
- Bollobás (JGT, 1990) showed that almost every graph has reconstruction number three, a **conjecture** by Harary and Plantholt (JGT, 1985).
- Kostochka, Nahvi, West and Zirlin (Eur. J Combin., 2021) proved that 3-regular graphs are 2-reconstructible.






- Wang obtained the following interesting result: Suppose that  $A$  and  $B$  are two integral symmetric matrices such that  $\det(xI - A) = \det(xI - B)$  and  $\det(xI - A_i) = \det(xI - B_i)$  for each  $i$ . If  $\det(xI - A)$  is irreducible over  $\mathbb{Q}[x]$ , then there exists a diagonal matrix  $D$  with each diagonal entry being  $\pm 1$  such that  $B = D^T A D$ .
- Although some reconstructible graphs are obtained, the Ulam's conjecture is still open.



Wei Wang, *A uniqueness theorem on matrices and reconstruction*, Journal of Combinatorial Theory, Ser. B, 99 (2009), 261–265.



## List of some recent references on the graph reconstruction

-  S. K. Gupta, P. Mangal, V. Paliwal, *Some work towards the proof of the reconstruction conjecture*, Discrete Mathematics, 272(2003), 291–296.
-  R. Forman, *Finite-type invariants for graphs and graph reconstructions*, Advances in Mathematics, 186 (2004), 181–228.
-  Wei Wang, *A uniqueness theorem on matrices and reconstruction*, Journal of Combinatorial Theory, Ser. B, 99 (2009), 261–265.
-  A. V. Kostochka, M. Nahvi, D. B. West and D. Zirlin, *3-regular graphs are 2-reconstructible*, European Journal of Combinatorics, 91(2021), 103216.
-  T. Hosaka, *The reconstruction conjecture for finite simple graphs and associated directed graphs*, Discrete Mathematics, 345 (2022), 112893.

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## Edge Reconstruction Conjecture of Graphs

- Harary posed a similar conjecture (i.e., the [edge reconstruction conjecture](#)), which states that every simple graph  $G$  with edge set  $E(G)$  can be reconstructed from its edge deck  $\{G - e | e \in E(G)\}$  for  $|E(G)| \geq 4$ .

Harary's Conjecture :  $\{G - e | e \in E(G)\} \Rightarrow G$  if  $|E(G)| \geq 4$ ?

- Lovász (JCTB, 1972) proved that the edge reconstruction conjecture holds for simple graphs with  $n$  vertices and at least  $\frac{n(n-1)}{4}$  edges.

If  $|E(G)| \geq n(n-1)/4$ , then  $\{G - e | e \in E(G)\} \Rightarrow G$ .



F. Harary, *On the reconstruction of a graph from a collection of subgraphs*, Theory of Graphs and Its Applications (M. Fiedler, ed.), Czechoslovak Academy of Sciences, Prague/Academic Press, New York, 1965, 47–52.

## Edge Reconstruction Conjecture of Graphs

- Müller (JCTB, 1977) improved the Lovász's result and proved that the edge reconstruction conjecture holds for simple graphs with  $n$  vertices and more than  $n \cdot \log_2 n$  edges.

If  $|E(G)| \geq n \log_2 n$ , then  $\{G - e | e \in E(G)\} \Rightarrow G$ .

- Godsil, Krasikov and Roditty (JCTB, 1987) proved that if  $2^{m-k} > n!$  or if  $2m > \binom{n}{2} + k$ , then  $G$  is reconstructible from its collection of  $k$ -edge deleted subgraphs, where  $m = |E(G)|$ .
- Although Müller's result implies that the edge reconstruction conjecture holds for almost all simple graphs, **the Harary's conjecture is still open**.
- If  $G$  is reconstructible and has no isolated vertices, then  $G$  is edge reconstructible, i.e.,

*Ulam's Conj. holds  $\Rightarrow$  Harary's Conj. holds.*

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# Reconstruction of Characteristic Polynomials

- Tutte proved that the characteristic polynomial  $f(G; x)$  of a graph  $G$  can be reconstructed from the deck  $\{G - v | v \in V\}$ .

$$\{G - v | v \in V\} \Rightarrow f(G; x).$$

- Tutte proved that if the characteristic polynomial  $f(G; x)$  of a graph  $G$  is irreducible over the rationals, then  $G$  is reconstructible.

$$\{G - v | v \in V\} \Rightarrow G \text{ if } f(G; x) \text{ is irreducible over } \mathcal{Q}.$$



W. T. Tutte, All the king's horses, in: Graph Theory and Related Topics, edited by J. A. Bondy and U. S. R. Murty (Academic Press, New York), 1979, 15–33.

# Reconstruction of Characteristic Polynomials

- Cvetković, at the XVIII International Scientific Colloquium in Ilmenau in 1973, posed a related problem as follows: Can the characteristic polynomial  $f(G; x)$  of a simple graph  $G$  with vertex set  $V(G)$  be reconstructed from  $\{f(G - v; x) \mid v \in V(G)\}$  for  $|V(G)| \geq 3$ ?

$$\{f(G - v; x) \mid v \in V(G)\} \Rightarrow f(G; x) \text{ if } |V(G)| \geq 3?$$

- The same problem was independently posed by Schwenk. Gutman and Cvetković obtained some results related to this problem.



A. J. Schwenk, *Spectral reconstruction problems*, Ann. New York Acad. Sci., 328 (1979), 183–189.



I. Gutman, D. Cvetković, *The reconstruction problem for characteristic polynomials of graphs*, Univ. Beograd, Publ. Elektrotehn. Fak. Ser. Mat. Fiz. No. 498 (451) (1975), 45–48.

## Reconstruction of Characteristic Polynomials

- Hagos proved that  $f(G; x)$  is reconstructible from  $\{f(G - v; x) \mid v \in V(G)\} \cup \{f((G - v)^c; x) \mid v \in V(G)\}$ .

$$\{f(G - v; x) \mid v \in V(G)\} \cup \{f((G - v)^c; x) \mid v \in V(G)\} \Rightarrow f(G; x).$$

- Note that  $f'(G; x) = \sum_{v \in V(G)} f(G - v; x)$ . Hence the coefficients of  $f(G; x)$  except for the constant term can be reconstructed from  $\{f(G - v; x) \mid v \in V(G)\}$ .
- No examples of non-unique reconstruction of the characteristic polynomial of graphs are known.



E. M. Hagos, *The characteristic polynomial of a graph is reconstructible from the characteristic polynomials of its vertex-deleted subgraphs and their complements*, the electronic journal of combinatorics, 7 (2000), #R12



- Under the assumption that the reconstruction of the characteristic polynomial is not unique, Cvetković described some properties of graphs  $G$  such that the constant term of  $f(G; x)$  can not be reconstructed from  $\{f(G - v; x) \mid v \in V(G)\}$ .
- Sciriha and Stanić survey classical and some more recent results concerning the reconstruction problem of the characteristic polynomial of graphs.



D. Cvetković, *On the reconstruction of the characteristic polynomial of a graph*, Discrete Mathematics, 212 (2000), 45–52.



I. Sciriha and Z. Stanić, *The polynomial reconstruction problem: The first 50 years*, Discrete Mathematics, 346 (2023), 113349.

- **A natural problem is: Can the characteristic polynomial  $f(G; x)$  (resp. permanental polynomial  $g(G; x)$ ) of a graph or digraph be reconstructed from  $\{f(G - e; x) | e \in E(G)\}$  (resp.  $\{g(G - e; x) | e \in E(G)\}$ )?**

$$\{f(G - e; x) | e \in E(G)\} \Rightarrow f(G; x) \text{ if } |E(G)| \geq 4?$$

$$\{g(G - e; x) | e \in E(G)\} \Rightarrow g(G; x) \text{ if } |E(G)| \geq 4?$$

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## Two Equalities on Determinants and Permanents

- Let  $X = (x_{st})_{n \times n}$  be a matrix of order  $n$  over the complex field. For any  $1 \leq i, j \leq n$ , define a matrix  $X_{ij} = (x_{st}^{ij})_{n \times n}$ , where

$$x_{st}^{ij} = \begin{cases} x_{st} & \text{if } (s, t) \neq (i, j) \\ 0 & \text{if } (s, t) = (i, j). \end{cases}$$

That is,  $X_{ij}$  is the matrix obtained from matrix  $X$  by replacing the  $(i, j)$ -entry  $x_{ij}$  with 0. For example, if

$$X = \begin{pmatrix} a & b & c \\ e & f & g \\ r & s & t \end{pmatrix},$$

then

$$X_{12} = \begin{pmatrix} a & 0 & c \\ e & f & g \\ r & s & t \end{pmatrix}, X_{33} = \begin{pmatrix} a & b & c \\ e & f & g \\ r & s & 0 \end{pmatrix}.$$

Obviously, if  $x_{ij} = 0$ , then  $X = X_{ij}$ .

## Two Equalities on Determinants and Permanents

- Let  $\varphi : S_n \rightarrow \mathbb{R}$  be a function, where  $S_n$  is the symmetric group of order  $n$ .
- The  $\varphi$ -immanant of  $X$  is defined as

$$\text{Imm}_\varphi(X) = \sum_{\alpha \in S_n} \varphi(\alpha) x_{1\alpha(1)} x_{2\alpha(2)} \cdots x_{n\alpha(n)},$$

where the sum ranges over all elements  $\alpha$  of  $S_n$ .

- Obviously,  $\text{Imm}_\varphi(X) = \det(X)$  if  $\varphi(\alpha) = \text{sgn}(\alpha)$  for any  $\alpha \in S_n$  and  $\text{Imm}_\varphi(X) = \text{per}(X)$  if  $\varphi(\beta) = 1$  for any  $\beta \in S_n$ .

### Theorem 1 (Zhang, Jin, Yan, 2023)

Let  $X = (x_{st})_{n \times n}$  be a matrix of order  $n$  over the complex field and let  $X_{ij}$  be defined as above. Then for any function  $\varphi : S_n \rightarrow \mathbb{R}$ , the  $\varphi$ -immanant of  $X$  satisfies:

$$(n^2 - n)\text{Imm}_\varphi(X) = \sum_{1 \leq i, j \leq n} \text{Imm}_\varphi(X_{ij}). \quad (1)$$

## Two Equalities on Determinants and Permanents

The following theorem is immediate from Theorem 1.

### Theorem 2 (Zhang, Jin, Yan, 2023)

Let  $X = (x_{st})_{n \times n}$  be a matrix of order  $n$  over the complex field and let  $X_{ij}$  be defined as above. Then the determinant of  $X$  satisfies:

$$(n^2 - n) \det(X) = \sum_{1 \leq i, j \leq n} \det(X_{ij}). \quad (2)$$

### Theorem 3 (Zhang, Jin, Yan, 2023)

Let  $X = (x_{st})_{n \times n}$  be a matrix of order  $n$  over the complex field and let  $m$  be the number of non-zero entries of  $X$ . Then

$$(m - n) \det(X) = \sum_{(i,j) \in I} \det(X_{ij}), \quad (3)$$

where  $I = \{(i, j) | x_{ij} \neq 0, 1 \leq i, j \leq n\}$ .



## Two Equalities on Determinants and Permanents

- By the definition of the determinant and permanent,

$$\det(X) = \sum_{\alpha \in S_n} \operatorname{sgn}(\alpha) x_{1,\alpha(1)} x_{2,\alpha(2)} \cdots x_{n,\alpha(n)},$$

$$\operatorname{per}(X) = \sum_{\alpha \in S_n} x_{1,\alpha(1)} x_{2,\alpha(2)} \cdots x_{n,\alpha(n)}.$$

### Theorem 4 (Zhang, Jin, Yan, 2023)

Let  $X = (x_{st})_{n \times n}$  be a matrix of order  $n$  over the complex field and let  $m$  be the number of non-zero entries of  $X$ . Then the permanent  $\operatorname{per}(X)$  of  $X$  satisfies:

$$(m - n)\operatorname{per}(X) = \sum_{(i,j) \in I} \operatorname{per}(X_{ij}), \quad (4)$$

where  $I = \{(i,j) | x_{ij} \neq 0, 1 \leq i, j \leq n\}$ .

# Proof of Theorem 1

- For convenience, we assume that  $\{x_{11}, x_{12}, \dots, x_{nn}\}$  is a set of  $n^2$  independent commuting variables over the complex field.
- There exist  $n!$  terms in the expansion of the  $\varphi$ -immanant  $\text{Imm}_\varphi(X)$  of  $X$ , denoted by  $\mathcal{T}_{\alpha_1}, \mathcal{T}_{\alpha_2}, \dots, \mathcal{T}_{\alpha_{n!}}$ , where  $S_n = \{\alpha_i | 1 \leq i \leq n\}$  is the symmetric group of order  $n$  and

$$\mathcal{T}_{\alpha_i} = \varphi(\alpha_i) x_{1\alpha_i(1)} x_{2\alpha_i(2)} \cdots x_{n\alpha_i(n)}.$$

- Hence

$$\text{Imm}_\varphi(X) = \sum_{i=1}^{n!} \mathcal{T}_{\alpha_i}.$$

- Define an  $n! \times n^2$  matrix  $T = (t_{\alpha_i X_{st}})_{n! \times n^2}$  as

$$t_{\alpha_i X_{st}} = \begin{cases} \mathcal{T}_{\alpha_i}, & \text{if } (s, t) \notin \{(k, \alpha_i(k)) | 1 \leq k \leq n\}; \\ 0, & \text{otherwise.} \end{cases}$$

## Proof of Theorem 1

- Hence each row of  $T$  has  $n^2 - n$  entries each of which equals  $\mathcal{T}_{\alpha_i}$  and  $n$  entries each of which equals zero.
- For example, if  $n = 3$ , then  $U = S_3 = \{\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5, \alpha_6\}$  and  $V = \{x_{11}, x_{12}, x_{13}, x_{21}, x_{22}, x_{23}, x_{31}, x_{32}, x_{33}\}$ , where

$$\alpha_1 = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}, \alpha_2 = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}, \alpha_3 = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix},$$

$$\alpha_4 = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}, \alpha_5 = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix}, \alpha_6 = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix}.$$

# Proof of Theorem 1

$$\mathcal{T}_{\alpha_1} = \varphi(\alpha_1)x_{11}x_{22}x_{33}, \mathcal{T}_{\alpha_2} = \varphi(\alpha_2)x_{11}x_{23}x_{32},$$

$$\mathcal{T}_{\alpha_3} = \varphi(\alpha_3)x_{13}x_{22}x_{31}, \mathcal{T}_{\alpha_4} = \varphi(\alpha_4)x_{12}x_{21}x_{33},$$

$$\mathcal{T}_{\alpha_5} = \varphi(\alpha_5)x_{12}x_{23}x_{31}, \mathcal{T}_{\alpha_6} = \varphi(\alpha_6)x_{13}x_{21}x_{32}.$$

$$T = \begin{matrix} & \alpha_1 & \alpha_2 & \alpha_3 & \alpha_4 & \alpha_5 & \alpha_6 \\ \begin{matrix} \alpha_1 \\ \alpha_2 \\ \alpha_3 \\ \alpha_4 \\ \alpha_5 \\ \alpha_6 \end{matrix} & \begin{pmatrix} x_{11} & x_{12} & x_{13} & x_{21} & x_{22} & x_{23} & x_{31} & x_{32} & x_{33} \\ 0 & \mathcal{T}_{\alpha_1} & \mathcal{T}_{\alpha_1} & \mathcal{T}_{\alpha_1} & 0 & \mathcal{T}_{\alpha_1} & \mathcal{T}_{\alpha_1} & \mathcal{T}_{\alpha_1} & 0 \\ 0 & \mathcal{T}_{\alpha_2} & \mathcal{T}_{\alpha_2} & \mathcal{T}_{\alpha_2} & \mathcal{T}_{\alpha_2} & 0 & \mathcal{T}_{\alpha_2} & 0 & \mathcal{T}_{\alpha_2} \\ \mathcal{T}_{\alpha_3} & \mathcal{T}_{\alpha_3} & 0 & \mathcal{T}_{\alpha_3} & 0 & \mathcal{T}_{\alpha_3} & 0 & \mathcal{T}_{\alpha_3} & \mathcal{T}_{\alpha_3} \\ \mathcal{T}_{\alpha_4} & 0 & \mathcal{T}_{\alpha_4} & 0 & \mathcal{T}_{\alpha_4} & \mathcal{T}_{\alpha_4} & \mathcal{T}_{\alpha_4} & \mathcal{T}_{\alpha_4} & 0 \\ \mathcal{T}_{\alpha_5} & 0 & \mathcal{T}_{\alpha_5} & \mathcal{T}_{\alpha_5} & \mathcal{T}_{\alpha_5} & 0 & 0 & \mathcal{T}_{\alpha_5} & \mathcal{T}_{\alpha_5} \\ \mathcal{T}_{\alpha_6} & \mathcal{T}_{\alpha_6} & 0 & 0 & \mathcal{T}_{\alpha_6} & \mathcal{T}_{\alpha_6} & \mathcal{T}_{\alpha_6} & 0 & \mathcal{T}_{\alpha_6} \end{pmatrix} \end{matrix}.$$

- Note that the  $i$ -th row in matrix  $T$  has exactly  $n^2 - n$  non-zero entries each of which equals  $\mathcal{T}_{\alpha_i}$ . Hence the sum of entries of  $i$ -th row in  $T$  equals  $(n^2 - n)\mathcal{T}_{\alpha_i}$ . So the sum of all entries in  $T$  equals  $(n^2 - n)\text{Imm}_\varphi(X)$ . That is,

$$\sum_{i=1}^{n!} \sum_{1 \leq s, t \leq n} t_{\alpha_i \times st} = (n^2 - n)\text{Imm}_\varphi(X). \quad (5)$$

- On the other hand, by the definition of  $X_{st}$ , the sum of entries of  $x_{st}$ -th column of  $T$  equals exactly the determinant of  $X_{st}$ . Hence

$$\sum_{1 \leq s, t \leq n} \sum_{i=1}^{n!} t_{\alpha_i x_{st}} = \sum_{1 \leq s, t \leq n} \text{Imm}_{\varphi}(X_{st}). \quad (6)$$

- Hence, by Eqs. (5) and (6), the theorem holds.

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## Six Digraph Polynomials

- Let  $G = (V(G), E(G))$  be a digraph having no loops and no multiple arcs, with vertex set  $V(G) = \{v_1, v_2, \dots, v_n\}$  and arc set  $E(G) = \{e_1, e_2, \dots, e_m\}$ .
- Denote the adjacency matrix and the vertex in-degree diagonal matrix of  $G$  by  $A = (a_{ij})_{n \times n}$  and  $D = \text{diag}(d^+(v_1), d^+(v_2), \dots, d^+(v_n))$ , where  $a_{ij} = 1$  if  $(v_i, v_j) \in E(G)$  and  $a_{ij} = 0$  otherwise, and  $d^+(v)$  is the number of arcs with head  $v_j$ .
- Then  $D - A$  and  $D + A$  are the Laplacian and signless Laplacian matrices of  $G$ .



- Set

$$f_1(G; x) = \det(xI - A),$$

$$f_2(G; x) = \det(xI - D + A),$$

$$f_3(G; x) = \det(xI - D - A),$$

$$f_4(G; x) = \text{per}(xI - A),$$

$$f_5(G; x) = \text{per}(xI - D + A),$$

$$f_6(G; x) = \text{per}(xI - D - A),$$

where  $\det(X)$  and  $\text{per}(X)$  denote the determinant and permanent of a square matrix  $X$ .

- Then  $f_1(G; x)$ ,  $f_2(G; x)$ ,  $f_3(G; x)$ ,  $f_4(G; x)$ ,  $f_5(G; x)$ , and  $f_6(G; x)$  are the characteristic polynomial, Laplacian characteristic polynomial, signless Laplacian characteristic polynomial, permanental polynomial, Laplacian permanental polynomial, signless Laplacian permanental polynomial of digraph  $G$ , respectively.

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# Main Results

- Suppose that both  $\beta$  and  $\gamma$  are real numbers satisfying  $\gamma \neq 0$ . Let  $G$  be a digraph with vertex set  $V(G) = \{v_1, v_2, \dots, v_n\}$  and arc set  $E(G) = \{e_1, e_2, \dots, e_m\}$ , which has no loops or multiple arcs. Set

$$g_1(G; x) = \det(xI_n - \beta D - \gamma A),$$

$$g_2(G; x) = \text{per}(xI_n - \beta D - \gamma A).$$

## Theorem 5 (Zhang, Jin, Yan, 2023)

Both  $g_1(G; x)$  and  $g_2(G; x)$  defined as above satisfy:

$$(m - n)g_1(G; x) + xg_1'(G; x) = \sum_{e \in E(G)} g_1(G - e; x), \quad (7)$$

$$(m - n)g_2(G; x) + xg_2'(G; x) = \sum_{e \in E(G)} g_2(G - e; x). \quad (8)$$

## Theorem 6 (Zhang, Jin, Yan, 2023)

Let  $G$  be a digraph with vertex set  $V(G) = \{v_1, v_2, \dots, v_n\}$  and arc set  $E(G) = \{e_1, e_2, \dots, e_m\}$ , which has no loops or multiple arcs. Then, for any  $i = 1, 2, 3, 4, 5, 6$ ,

$$(m - n)f_i(G; x) + xf'_i(G; x) = \sum_{e \in E(G)} f_i(G - e; x). \quad (9)$$

## Theorem 7 (Zhang, Jin, Yan, 2023)

Let  $G$  be a digraph with vertex set  $V(G) = \{v_1, v_2, \dots, v_n\}$  and arc set  $E(G) = \{e_1, e_2, \dots, e_m\}$ , which has no loops or multiple arcs. If  $m \neq n$ , then  $f_i(G; x)$  can be reconstructed from  $\{f_i(G - e; x) | e \in E(G)\}$  for  $1 \leq i \leq 6$ .

$$\{f_i(G - e; x) | e \in E(G)\} \Rightarrow f_i(G; x) \text{ for } 1 \leq i \leq 6 \text{ if } m \neq n.$$

# Main Results

- Let  $G_1 = C_n$  be a directed cycle with vertex set  $V(G_1) = \{v_i | 1 \leq i \leq n\}$  and arc set  $E(G_1) = \{(v_i, v_{i+1}) | 1 \leq i \leq n-1\} \cup \{(v_n, v_1)\}$ .
- Let  $G_2$  be the digraph obtained from  $G_1$  by replacing arc  $(v_n, v_1)$  with  $(v_1, v_n)$ .

$$\{f_i(G_1 - e; x) | e \in E(G_1)\} = \{f_i(G_2 - e) | e \in E(G_2)\},$$

but  $f_i(G_1; x) \neq f_i(G_2; x)$  for  $i = 1, 4$ .

## Theorem 8 (Zhang, Jin, Yan, 2023)

*If  $m = n$ , the characteristic polynomial  $f_1(G; x)$  and permanent polynomial  $f_4(G; x)$  of a digraph  $G$  with  $n$  vertices and  $n$  arcs can not be determined uniquely by  $\{f_1(G - e; x) | e \in E(G_1)\}$  and  $\{f_4(G - e; x) | e \in E(G_1)\}$ , respectively. That is,  $\{f_i(G - e; x) | e \in E(G)\} \not\Rightarrow f_i(G; x)$  for  $i = 1, 4$ , if  $m = n$ .*

- Note that,  $f_2(G; x) = \det(xI - D + A)$ . Hence  $f_2(G; 0) = \det(A - D) = 0$ . If  $m = n$ , then

$$xf_2'(G; x) = \sum_{e \in E(G)} f_2(G - e; x).$$

Given the initial condition  $f_2(G; 0) = 0$ , the differential equation above has a unique solution.

## Theorem 9 (Zhang, Jin, Yan, 2023)

*The Laplacian characteristic polynomial  $f_2(G; x) = \det(xI - D + A)$  of a digraph  $G$  with arc set  $E(G)$  can be reconstructed from  $\{f_2(G - e; x) | e \in E(G)\}$ . That is,  $\{f_2(G - e; x) | e \in E\} \Rightarrow f_2(G; x)$ .*

## 1 Introduction

- Vertex Reconstruction Conjecture of Graphs
- Edge Reconstruction Conjecture of Graphs
- Reconstruction of Characteristic Polynomials

## 2 Edge Reconstruction of Digraph Polynomials

- Two Equalities on Determinants and Permanents
- Six Digraph Polynomials
- Main Results

## 3 Discussion

- For a digraph  $G = (V(G), E(G))$ ,

$$\{f_i(G - e; x) | e \in E(G)\} \Rightarrow f_i(G; x) \text{ for } 1 \leq i \leq 6, \text{ if } m \neq n.$$

$$\{f_i(G - e; x) | e \in E(G)\} \not\Rightarrow f_i(G; x) \text{ for } i = 1, 4, \text{ if } m = n.$$

$$\{f_2(G - e; x) | e \in E(G)\} \Rightarrow f_2(G; x), \text{ if } m = n.$$

$$\{f_i(G - e; x) | e \in E(G)\} \Rightarrow f_i(G; x) \text{ for } i = 3, 5, 6, \text{ if } m = n?$$



- **A natural problem is: Do the above results hold for simple graphs?**
- Let  $G$  be a simple graph with adjacency matrix  $A$  and vertex degree diagonal matrix  $D$ . Set

$$h_1(G; x) = \det(xI - A),$$

$$h_2(G; x) = \det(xI - D + A),$$

$$h_3(G; x) = \det(xI - D - A),$$

$$h_4(G; x) = \text{per}(xI - A),$$

$$h_5(G; x) = \text{per}(xI - D + A),$$

$$h_6(G; x) = \text{per}(xI - D - A).$$

## Theorem 10 (Zhang, Jin, Yan, Liu, 2023)

$$(m-n)h_1(G; x) + xh_1'(G; x) = \sum_{uv \in E(G)} [h_1(G - uv; x) + h_1(G - u - v; x)], \quad (10)$$

$$(m-n)h_4(G; x) + xh_4'(G; x) = \sum_{uv \in E(G)} [h_4(G - uv; x) - h_4(G - u - v; x)], \quad (11)$$

and for  $j = 2, 3$ ,

$$(m-n)h_j(G; x) + xh_j'(G; x) = \sum_{uv \in E(G)} h_j(G - uv; x). \quad (12)$$

- For a graph  $G = (V, E)$ , if  $m \neq n$ , then

$$\{h_i(G-e; x) | e \in E\} \cup \{h_i(G-u-v; x) | uv \in E\} \Rightarrow h_i(G; x) \text{ if } i = 1, 4,$$

$$\{h_i(G-e; x) | e \in E\} \Rightarrow h_i(G; x) \text{ for } i = 2, 3.$$

**Thank you for your attention!**