

Last time

- Graph container lemma (KW Alg)
- Count $\# C_4$ -free graphs on $[n] = \{1, 2, \dots, n\}$

$$\underline{ex(n, C_4)} = \Theta(n^{3/2})$$

$$f_n(C_4) \leq 2^{c \cdot n^{3/2}}$$

Application 1

Def A set $S \subseteq [n]$ is multiplicative

Sidon if \nexists distinct $a, b, c, d \in S$
satisfying $a \cdot b = c \cdot d$

ex $\{2, 3, 4, 6, 11, 20\}$

NOT multi. Sidon

$$2 \cdot 6 = 3 \cdot 4$$

ex Take {primes $p : p \leq n$ }

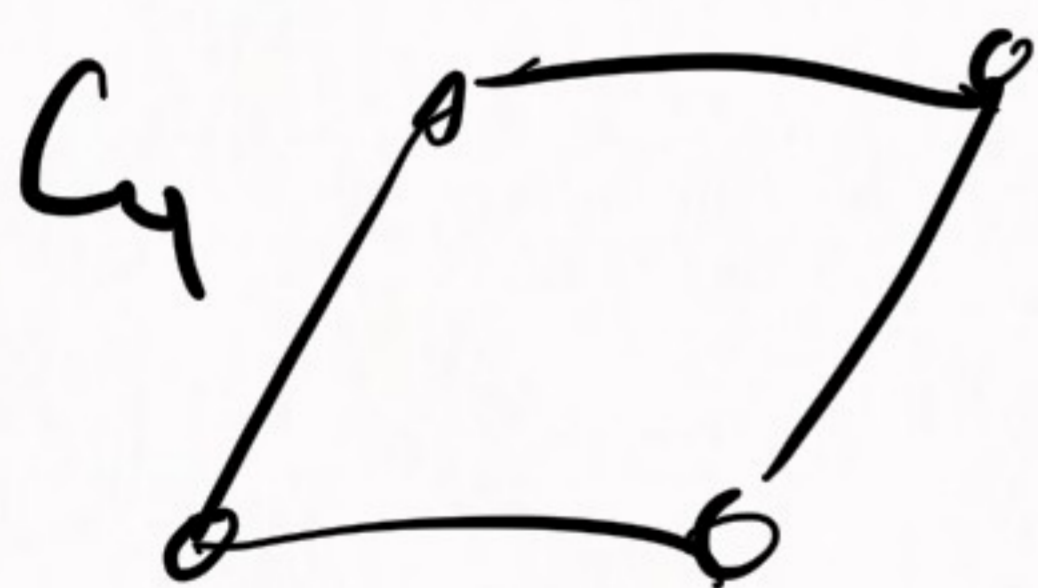
is multi. Sidon

No $ab = cd$

Erdős 38/69 $\left\{ \begin{array}{l} s(n) = \max \text{ size } S \subseteq [n] \\ S \text{ is multi. Sidon} \end{array} \right.$

$$s(n) = \pi(n) + \Theta\left(\frac{n^{3/4}}{\log^{3/2} n}\right) \leftrightarrow L$$

Lower bound $\left\{ \begin{array}{l} \text{primes btw } [n^{2/3}, n] \\ \text{auxiliary } G, V(G) = \text{primes} \\ \leq \sqrt{n} \\ \text{max-size } C_4\text{-free graph} \end{array} \right.$



$$L = \{u, v : uv \in E(G)\}$$

$$|G| = \frac{c\sqrt{n}}{\log n} \Rightarrow e(G) = \frac{c' n^{3/4}}{\log^{3/2} n}$$

Thm L. - Pach 19

$S(n) = \#$ multi-Sidon sets in $[n]$

$$S(n) = T(n) \cdot 2^{\Theta\left(\frac{n^{3/4}}{\log^{3/2} n}\right)}$$

Rank

$$1) T(n) = \prod_{\substack{n^{2/3} < p \leq n \\ \text{Primes}}} \left(\left\lfloor \frac{n}{p} \right\rfloor + 1 \right)$$

$$= 2^{\alpha + o(1)} \prod(n), \text{ where}$$

$$\alpha = \sum_{i=1}^{\infty} \frac{1}{i} \log_2 \left(1 + \frac{1}{i} \right) \approx \underline{1.8146} > 1$$

2) Let result up to l.o.t.

$$ex(n, m, C_4) = O(mn^{1/2} + m)$$

3) graph theory works when estimating

the l.o.t.

bipartite

Lem (Count $\#$ unbalanced $\wedge C_4$ -free graphs)

$$\forall \frac{1}{2} + \epsilon \leq m \leq n$$



$$\#(m, n)\text{-unbip. } C_4\text{-free graphs} = 2^{O(mn^{1/2})}$$

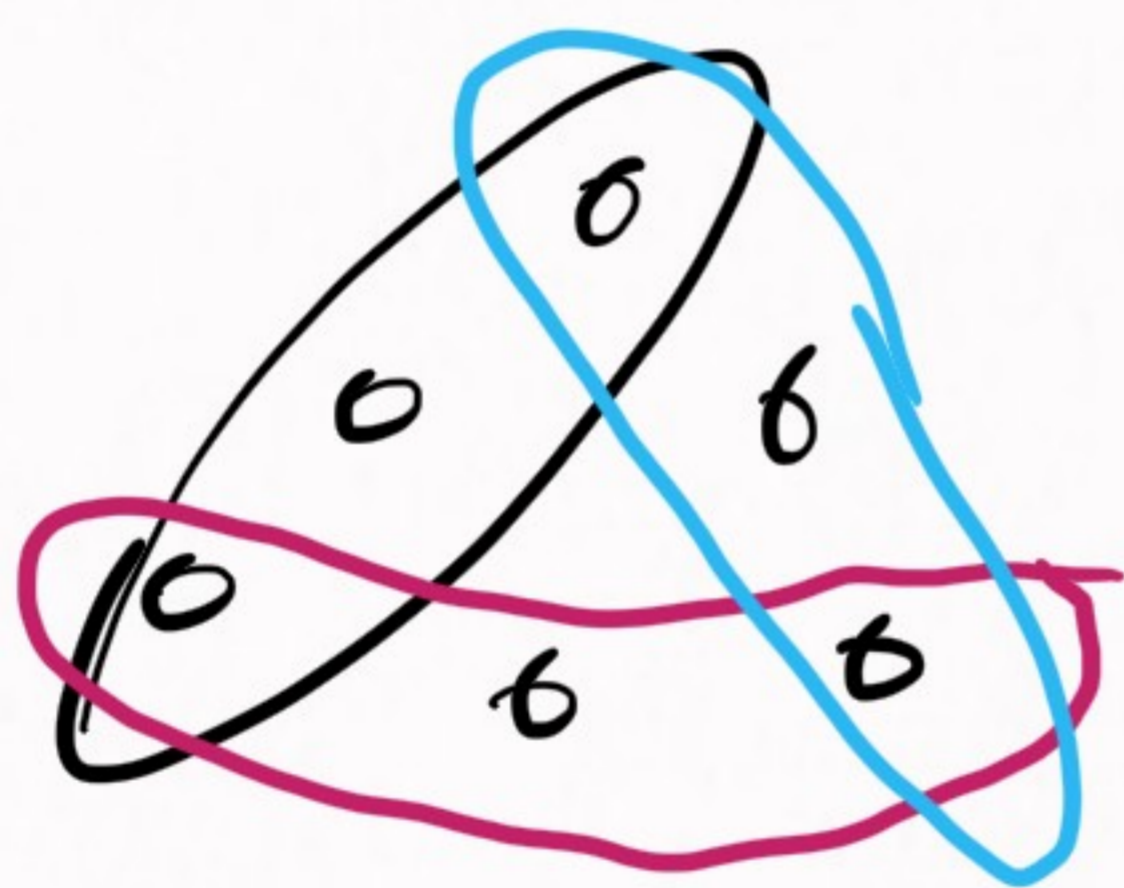
Perhaps

$\forall m \leq n$

(m, n) - \times bip. C_4 -free = $2^{O(mn^{1/2} + m)}$
graphs

Application 2

Def k -unif hypergraph \mathcal{H} is intersecting if
 $\forall A, B \in E(\mathcal{H}), A \cap B \neq \emptyset$



Erdős-Ko-Rado (1)

$\forall n \geq 2k,$ k -unif intersecting \mathcal{H}

$$\Rightarrow e(\mathcal{H}) \leq \binom{n-1}{k-1}$$

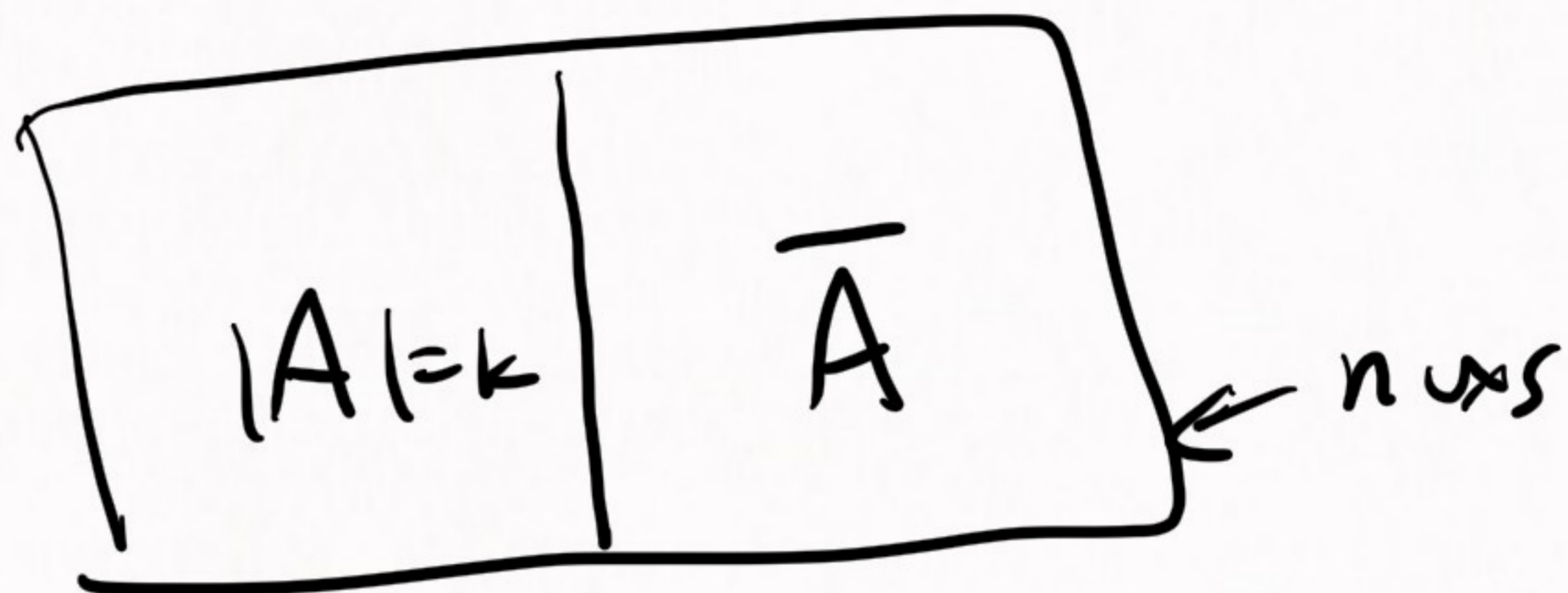
Thm Babai - Das - Delcourt - L. - Sharifzadeh 2015

$\forall n \geq 2k+1,$ # k -unif intersecting = $2^{(1+o(1)) \binom{n-1}{k-1}}$

Rank $n \geq 2k+1$ tight.

Consider $n=2k \Rightarrow \binom{n-1}{k-1} = \frac{1}{2} \binom{n}{k}$

$$\frac{\#(A, \bar{A})}{\frac{1}{2} \binom{n}{k}}$$



For each (A, \bar{A}) pair

Choose $\left\{ \begin{array}{l} \text{nothing} \\ A \\ \bar{A} \end{array} \right. \Rightarrow 3 = 3 \frac{1}{2} \binom{n}{k} \binom{n-1}{k-1}$

Goal Count k -unif intersecting hypergraphs

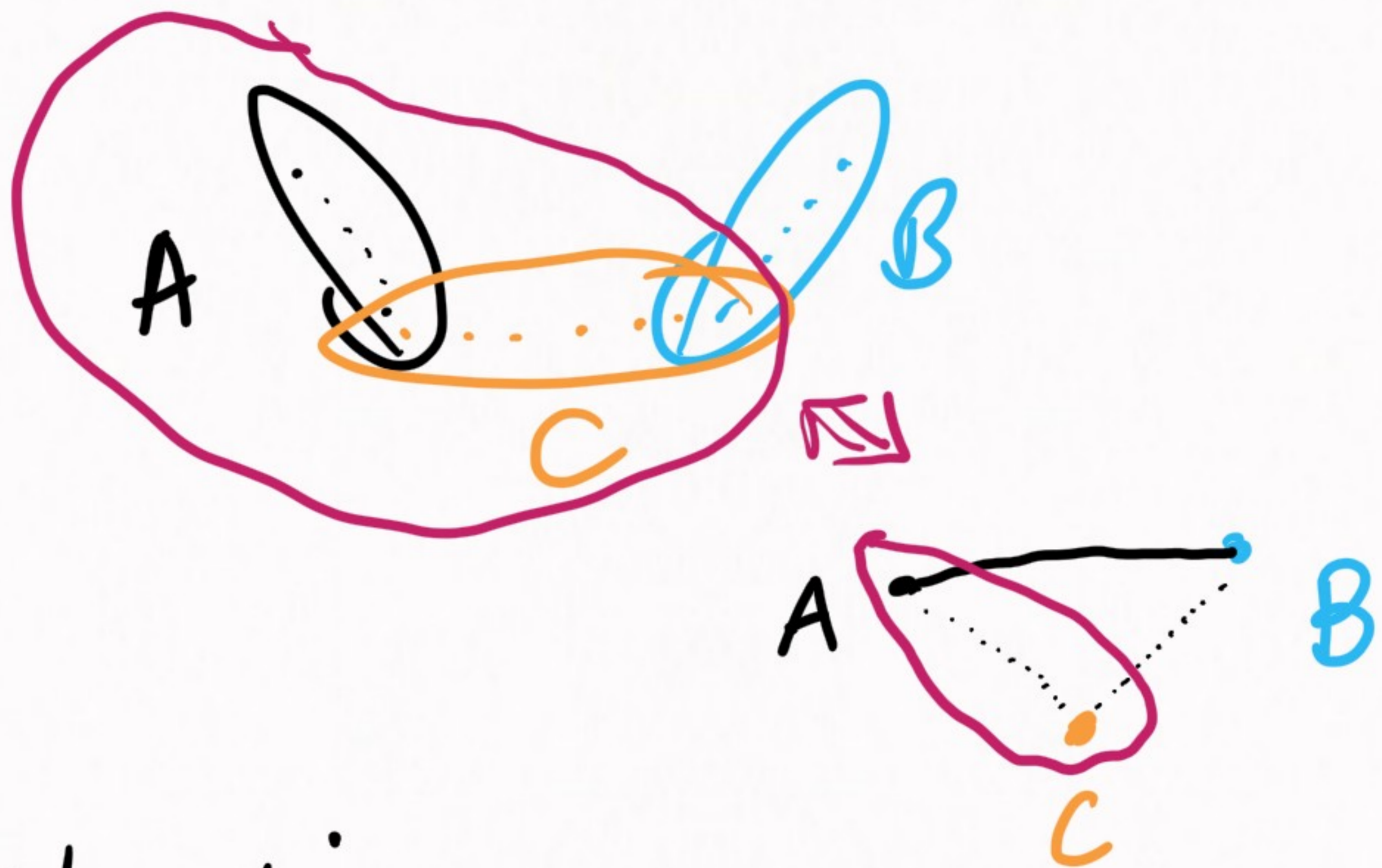
Count indep sets in certain graph

Kneser graph

Def Kneser graph $KG(n, k)$

$$V = \binom{[n]}{k}$$

$$E: A, B, \quad A \sim B \Leftrightarrow A \cap B = \emptyset$$



k -unif intersecting
hypergraph \mathcal{H}



an indep set
in Kneser graph.

Lem (Graph container)

• n -vx G

• R, β, ρ

$$R \geq n \cdot e^{-\beta \rho}$$



• locally dense

$$\forall u \subseteq V(G), |u| \geq R \Rightarrow e(G[u]) \geq \beta \binom{|u|}{2}$$

$$\exists \mathcal{C} \subseteq 2^{V(G)} \text{ st.}$$

$$|\mathcal{C}| \leq \binom{n}{\rho}$$

$$\forall C \in \mathcal{C}, |C| \leq \rho + R$$

• \forall indep set I

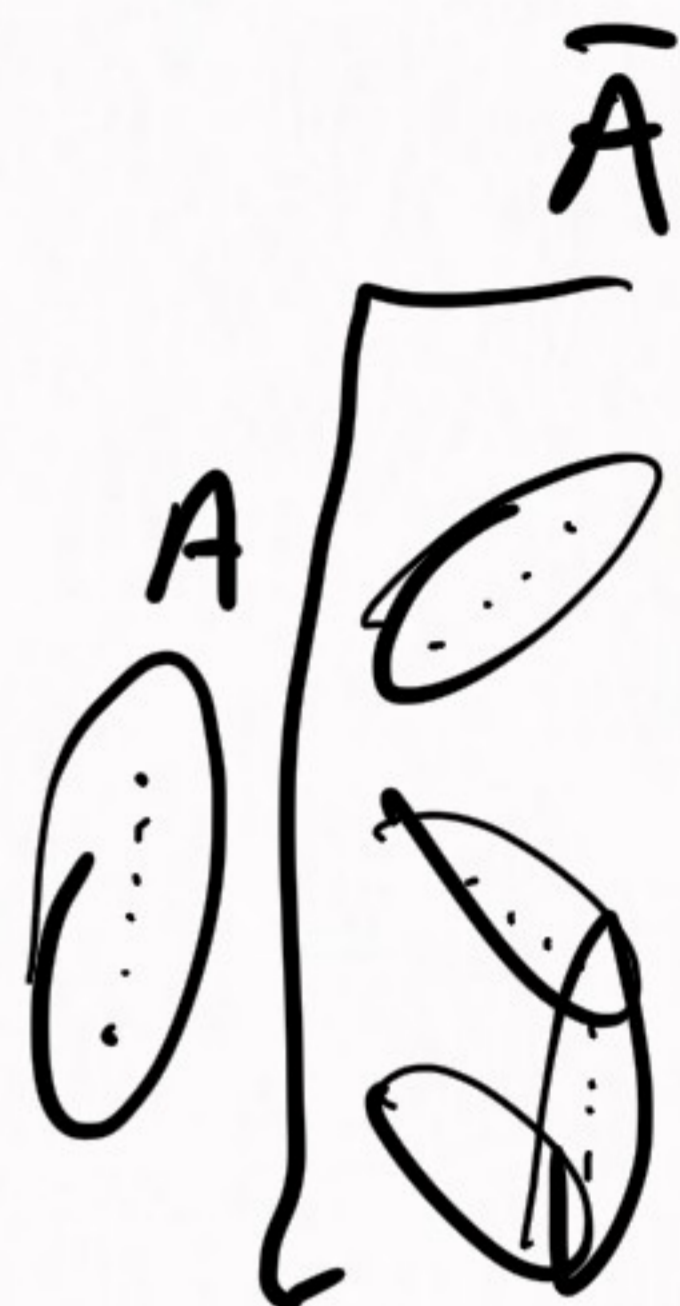
$$\exists C \in \mathcal{C} \text{ st. } I \subseteq C$$



Lovász 79 $\lambda_{\min} = \lambda = -\binom{n-k-1}{k-1}$

Kneser $N = \binom{n}{k}$

D -reg, $D = \binom{n-k}{k}$



Def (n, d, λ) -graph (pseudorandom)

n -vx, d -reg, $\max\{|\lambda_2|, |\lambda_n|\} \leq \lambda$

$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$

$\lambda \ll d \Rightarrow$ pseudorandom

Lem (Expander mixing Lem Alon-Chung 90s)

G n -vx

$\forall A \subseteq V(G)$

edge-density

d -reg \Rightarrow

$e(G[A]) \geq \frac{d}{n} \cdot \frac{|A|^2}{2} +$

$\lambda = \lambda_n$

expected see in random graph

$\frac{\lambda}{2n} |A| (n - |A|)$
error

Prop (supersaturation)

R

$\forall \epsilon > 0, \forall |S| \geq (1+\epsilon) \binom{n-1}{k-1}$ was in Kneser
graph $\Rightarrow e(G[S]) \geq \underbrace{\left(1 - \frac{1}{1+\epsilon}\right) \frac{D_n}{N(n-k)}}_{\text{blue underline}} \binom{|S|}{2}$

Remark 1. R almost ^{β} best we can hope for, as $R > \alpha(G) = \binom{n-1}{k-1}$
has to be

Application #3

Lower bound constr. in multicolour Ramsey.

Def $R(K_s, K_t) = \min n$ s.t.

\forall 2-edge-colouring of K_n \Rightarrow $\left\{ \begin{array}{l} K_s \\ \text{or} \\ K_t \end{array} \right.$
 k colours

$R(\triangle, \square, \text{pentagon})$

Recall $R(\Delta, K_m) = \underbrace{\Theta\left(\frac{m^2}{\log m}\right)}$

Open Const. $\left[\frac{1}{4} + o(1), 1 + o(1) \right]$

• Erdős-Sós 79 Conj

$$\lim_{m \rightarrow \infty} \frac{R(\Delta, \Delta, K_m)}{R(\Delta, K_m)} = \infty$$

That is, need a graph on $\Rightarrow \frac{m^2}{\log m}$ (construction)

s.t. $\begin{cases} \alpha(G) < m \\ G \text{ is } (\Delta, \Delta) \text{ free} \end{cases} \exists c_1, c_2 \text{ s.t.}$

Thm (Alon-Rödl 05) $m^3 \log^{c_2} m \geq \dots \geq m^3 \log^{c_1} m$

$R(\Delta, \Delta, K_m) = \Theta(m^3 \cdot \text{polylog } m)$

$R(\Delta, \dots, \Delta, K_m) = \Theta(m^{k+1} \text{ polylog } m)$

Goal: To show $R(\Delta, \Delta, K_m) > m^3 \cdot \text{polylog} m$

Construct G {

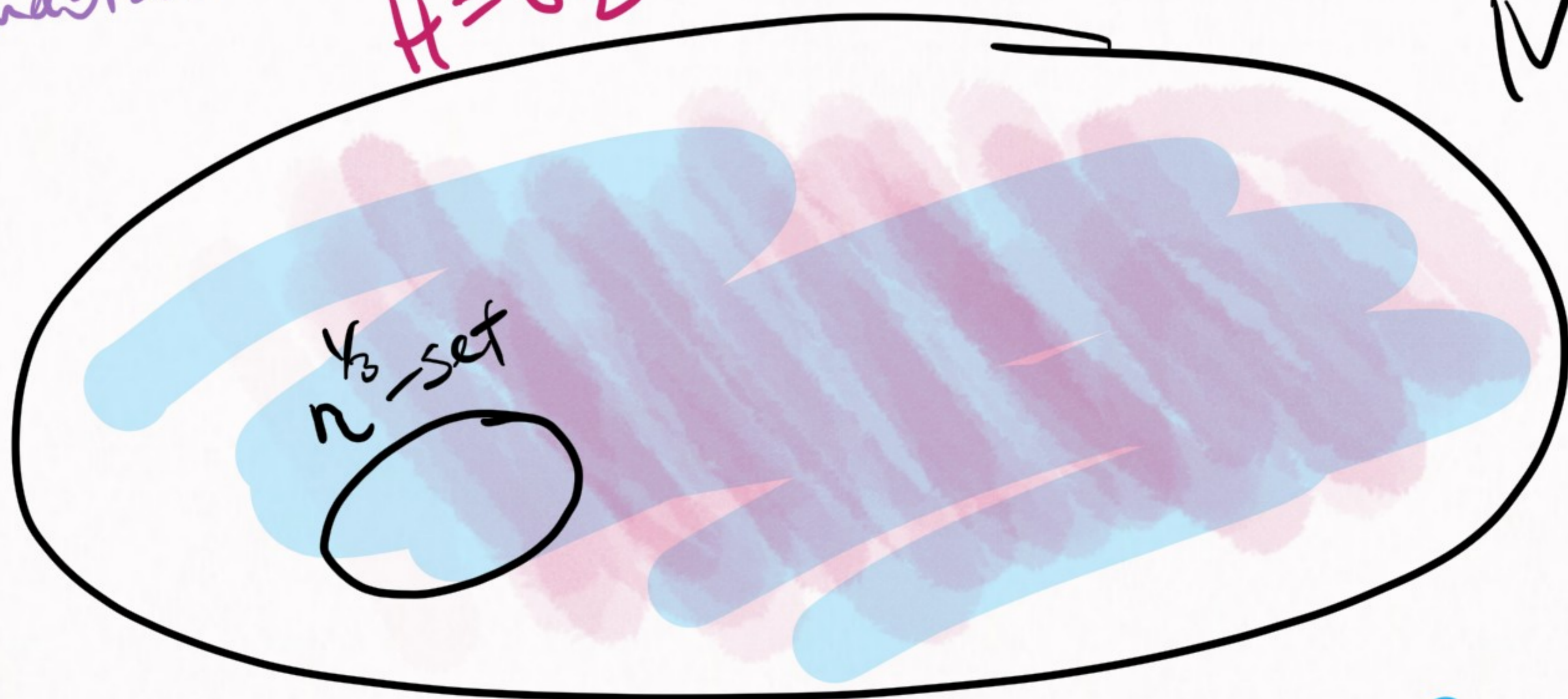
- (Δ, Δ) -free $E(G) = E(G_1) \cup E(G_2)$
- $\alpha(G) < m \approx n^{1/3} \text{polylog} n$ ↑ Δ -free
- $|G| = n \approx m^3 \text{polylog} m$.

PF Idea

Take H } $n \approx m$
} Δ -free
} "few" indep sets $\geq n^{1/3}$
} $H = G_2$

Let $G = G_1 \cup G_2$

such \nearrow
 H is dense pseudorandom Δ -free



Random $V(H) \rightarrow V$

$H = G_1$