

Non-bipartite k -common graphs

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Joint work with D. Král', J. Noel, S. Norin and F. Wei.

Ramsey multiplicities / Common graphs

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YES, if H is

Cycles, even-wheels, 5-wheel

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Fox-Wei ('17): H locally Sidorenko \equiv the girth of H is even

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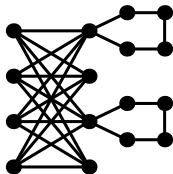
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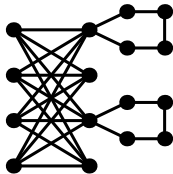
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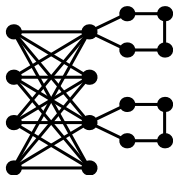
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Corollary: H locally Sidorenko $\iff H$ locally k -common $\forall k$

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$\#K_{\ell,\ell} \vee C_5 \geq d(G)^{\ell^2+5}$ by locally Sidorenko result of Fox-Wei?

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- 2) To every color class apply Lemma F or Lemma C ...

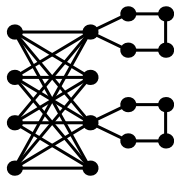
Lemma F: G is far from (pseudo)random $\implies \#H \geq d(G)^{e(H)}$
far $\equiv \#C_4 > d(G)^4 + \varepsilon \implies \#K_{\ell,\ell} \geq (d(G) + \frac{\varepsilon}{10})^{\ell^2}$ & $\#C_5 \gg 0$

Lemma C: G is close to (pseudo)random $\implies \#H \geq d(G)^{e(H)}$
 $\#K_{\ell,\ell} \vee C_5 \geq d(G)^{\ell^2+5}$ by locally Sidorenko result of Fox-Wei?

Proposition: Alternative spectral-based proof for the cases
 $K_{\ell,\ell} \vee C_5$ and $H_{2\ell,2\ell,C_5}$ that has no $\|\cdot\|_\infty$ assumption

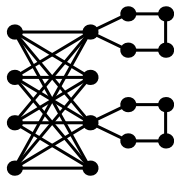
Conclusion

Theorem: $\forall k \geq 3 : \exists H_k$ with $\chi(H_k) = 3$ s.t. H_k is k -common



Conclusion

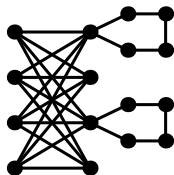
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Question: $\forall k, r \geq 3 : \exists H_k$ with $\chi(H_k) = r$ s.t. H_k is k -common

Conclusion

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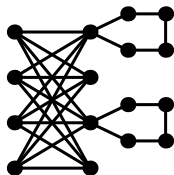


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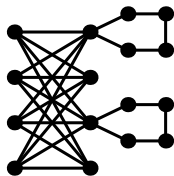
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Theorem: Fix $k \geq 3$. H locally k -common \iff girth of H is even

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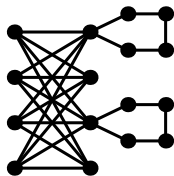
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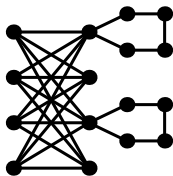
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Question: What are locally 2-common graphs?

YES: Sidorenko graphs, common graphs NO: K_4 (Cs3ka-Hubai-Lov3sz '19)

Conclusion Thank you for your attention!

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